Load Balancing for Interdependent IoT Microservices

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Outlines

Background and Motivation

System Modeling

Algorithm Design and Analysis

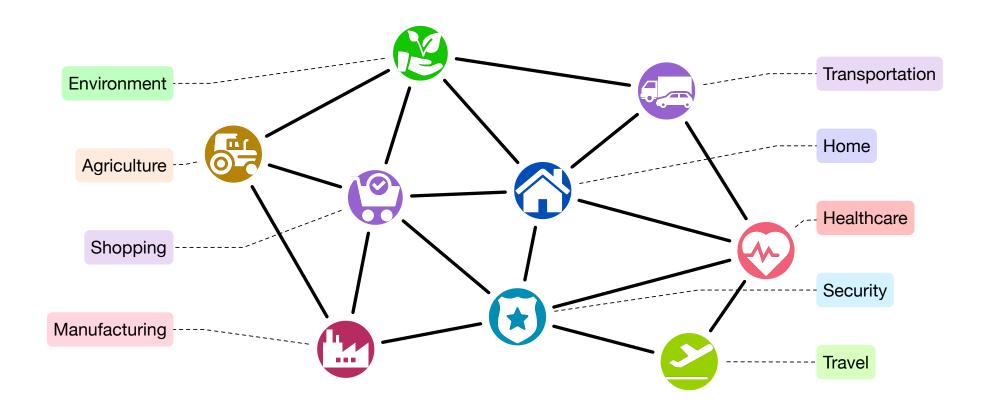
Performance Evaluation

Discussions, Future Work and Conclusions



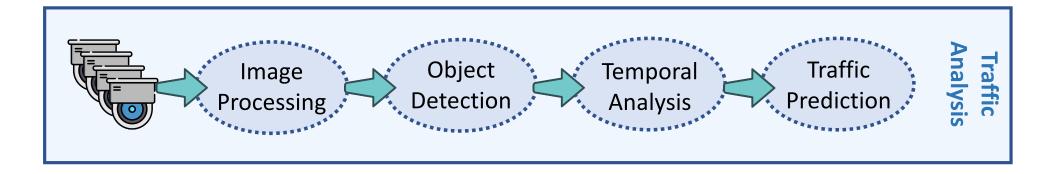
IoT:The Future Internet

■ IoT is the future Internet that connects every aspect of our work and life.



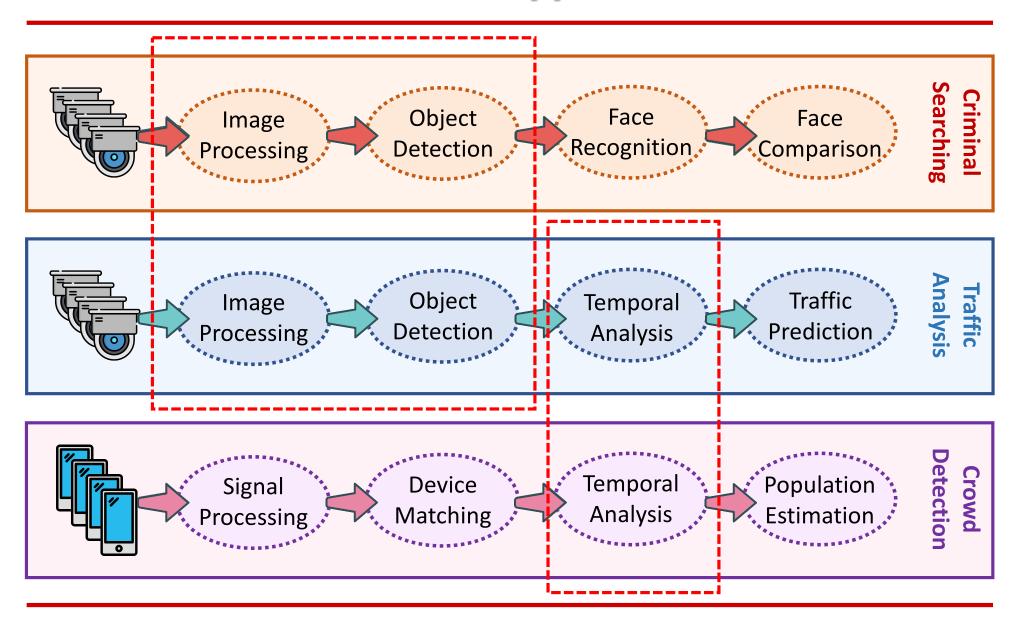


Example IoT Applications



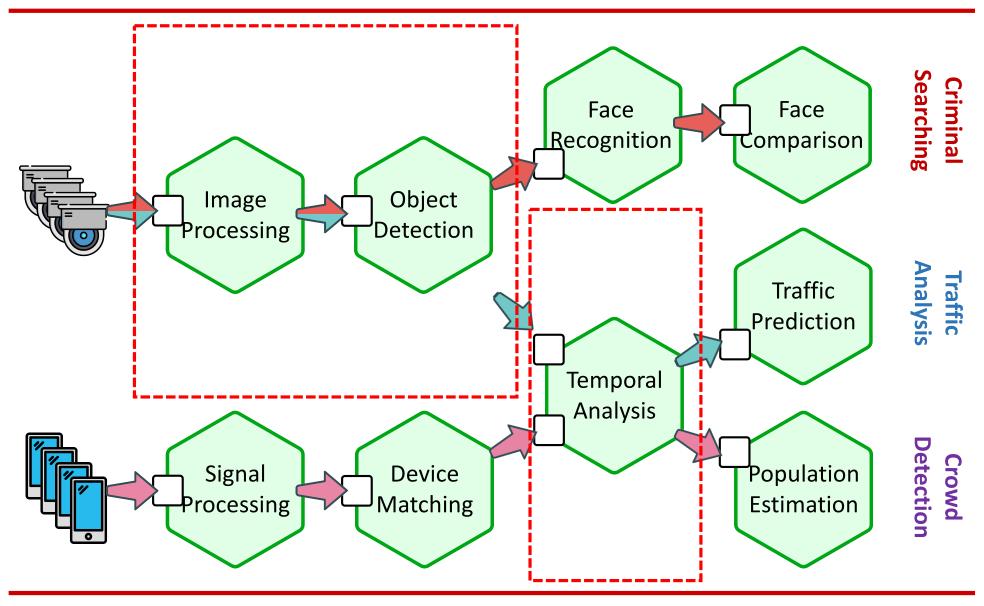


Monolithic Applications



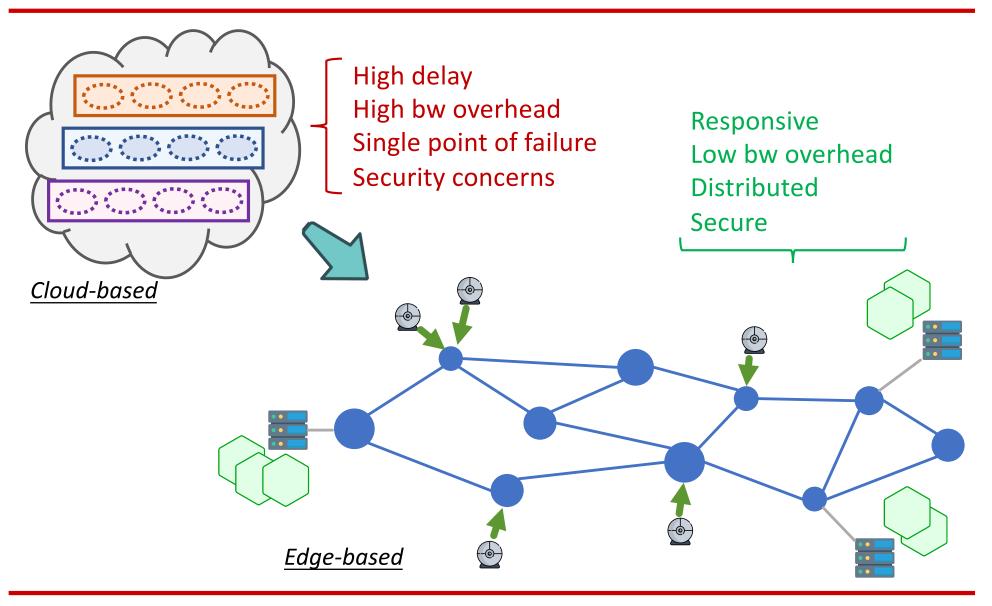


Microservices





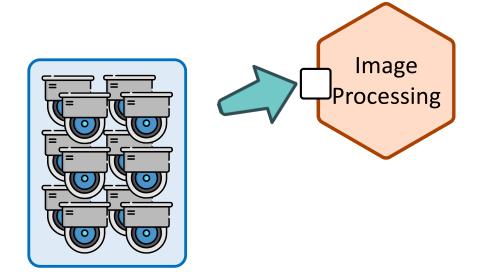
Microservices vs. Edge Computing





The Microservice Load Balancing Problem

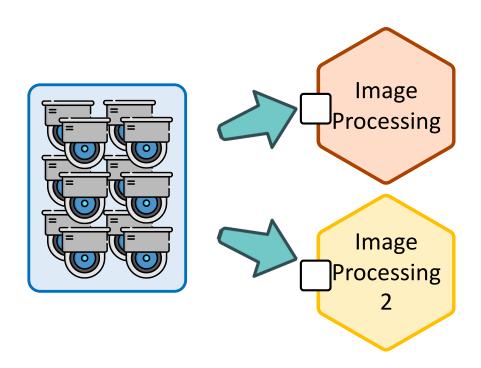
Edge-based microservices can be easily saturated.





The Microservice Load Balancing Problem

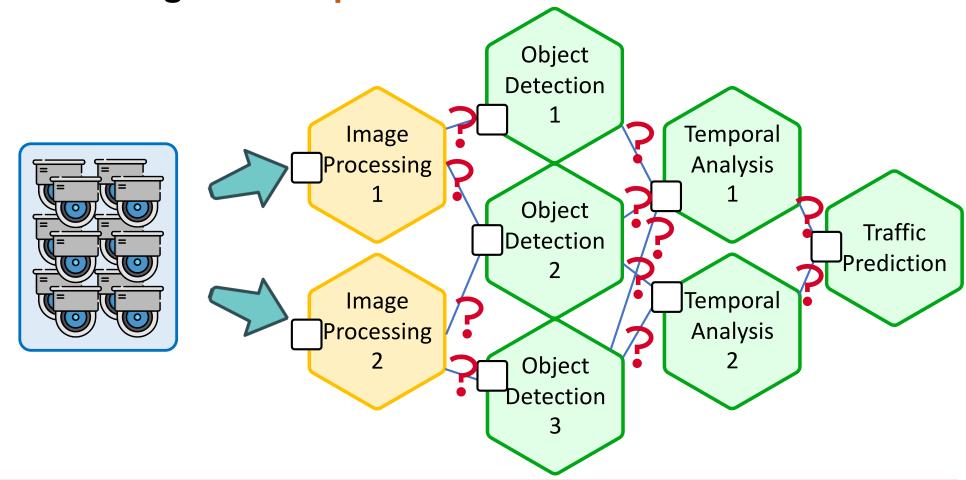
Edge-based microservices can be easily saturated.





The Microservice Load Balancing Problem

- Edge-based microservices can be easily saturated.
- Challenge: interdependent microservices.





Our Approach: Overview

Problem Modeling

- I) DAG-based interdependency graph (App-Graph).
- 2) Compactly modeled infrastructure (Inf-Graph).
- 3) Flexible application instantiation (Real-Graph).
- 4) Joint instantiation finding & load allocation.
- 5) Application QoS requirements.

Algorithmic Results

- 1) Optimal algorithm for QoS-agnostic problem.
- 2) NP-hardness for QoS-aware problem.
- 3) FPTAS for QoS-aware problem.

Next Steps (Future Work)

- Network-aware load balancing.
- 2) Reliability and security.
- 3) Economics-aware microservice composition.



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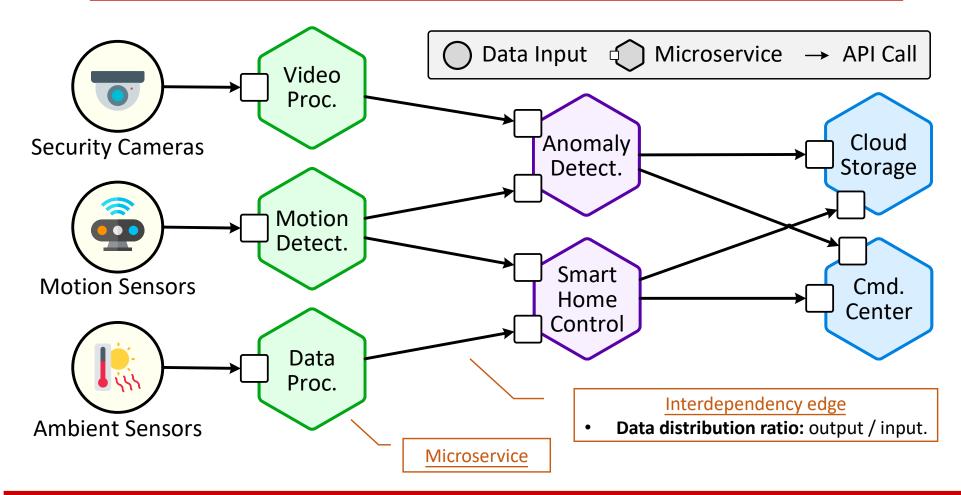
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Application with Interdependent Microservices

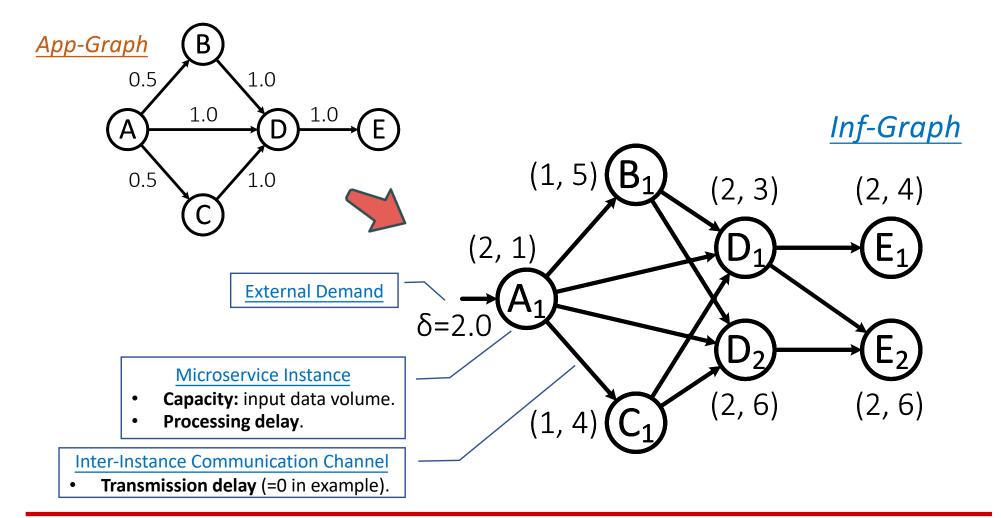
- General DAG-based application graph (App-Graph).
 - **Captures complex interdependencies, unlike existing line graph-based models.**





IoT Infrastructure in the Application's View

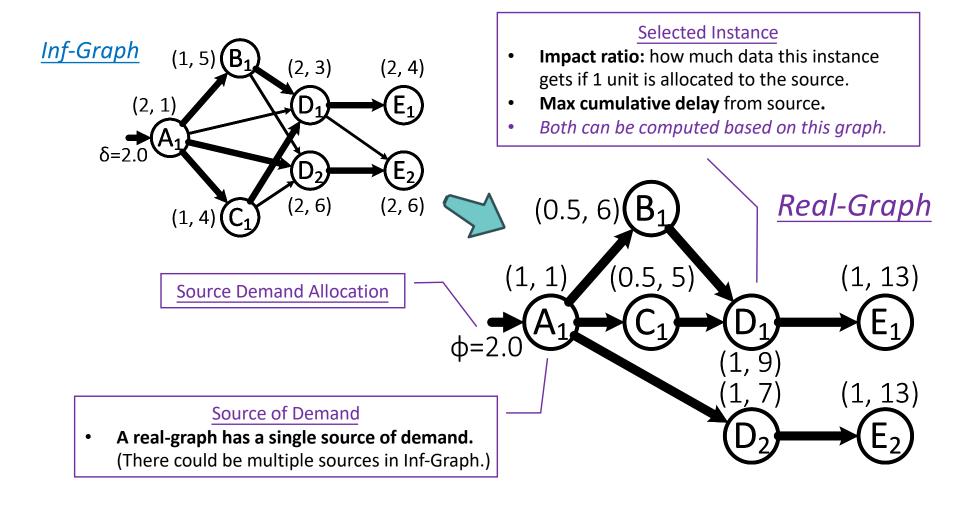
Inf-Graph: deployed microservice instances & their interactions.





Application Instantiation

Real-Graph: instantiating the App-Graph in the Inf-Graph.





Problem Statement: Overview

- ☐ Inputs:
 - App-Graph: microservices, interdependencies, data distribution ratios
 - ❖ Inf-Graph: instances, communication channels, capacities, and delays
- Outputs:
 - ❖ A set of Real-Graphs.
 - * External demand allocation for each Real-Graph.
- Constraints:
 - **\Load balancing**: total load \leq instance capacity * Ψ .
 - ❖ QoS awareness: maximum delay <= D.
- Objective (optimization version):
 - Minimize maximum delay of all Real-Graphs.



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BLB: Basic (QoS-agnostic) Load Balancing

Without delay constraint, the problem can be formulated as LP.

find $f: L \mapsto \mathbb{R}^*$

Variables: Per-link demand allocation function.

$$\delta_n \leq \Psi \cdot c_n, \quad \forall n \in N;$$

$$r_{(v_n,w)}\delta_n = \sum_{l \in L_{\text{out}}(n,w)} f(l), \quad \forall n, w \in V_{\text{out}}(v_n).$$

Capacity (load balancing) per node.

Flow conservation: sum flow towards all instances of a downstream microservice w = input data * data distribution ratio of w.

BLB: Basic (QoS-agnostic) Load Balancing

Without delay constraint, the problem can be formulated as LP.

 $\operatorname{find} \quad f:L\mapsto \mathbb{R}^*$ Variables: Per-link demand allocation function.

s.t.
$$\delta_n = \delta_n^{\mathrm{ext}} + \sum_{l \in L_{\mathrm{in}}(n)} f(l), \quad \forall n \in N;$$

Demand per node = external + flow-in.

$$\delta_n \leq \Psi \cdot c_n, \quad \forall n \in N;$$

$$r_{(v_n,w)}\delta_n = \sum_{l \in L_{\text{out}}(n,w)} f(l), \quad \forall n, w \in V_{\text{out}}(v_n).$$

Theorem I:

BLB is optimally solvable in polynomial time.

Real-Graph Decomposition Theorem

Theorem 2:

Every demand allocation function f so defined can be decomposed into at most |N| + |L| real-graphs with positive source demands.

- Why do we need such a theorem?
 - I. Transform any solution of BLB into a set of implementable real-graphs.
 - 2. Define QoS of a load balancing plan (max delay of all real-graphs).



QLB: QoS-aware Load Balancing

Theorem 3:

QLB (optimization version) is NP-hard.

☐ Fully Polynomial-Time Approximation Scheme (FPTAS)

can achieve the best trade-off between time and accuracy

- \clubsuit Approximation ratio: $(1+\epsilon)$ For maximization problem
- \bullet Time complexity: $O(\text{poly}(1/\epsilon) \times \text{poly}(\text{input}))$
- \clubsuit In practice, one can arbitrarily tune ϵ to get best accuracy within time limit.

Theorems 4&5:

QLB (optimization version) admits an FPTAS.



A Brief Overview of Our FPTAS

Idea:

- Pseudo-polynomial time algorithm:
 - Expand Inf-Graph into a delay-layered graph.
 - Run BLB LP on the expanded graph.
- Discretization via approximate testing:
 - Find delay lower & upper bounds (UB, LB) s.t. UB <= poly(input) * LB.
 - Discretize delay values based on (UB, LB).
 - Run pseudo-polynomial time algorithm.
 - Refine (UB, LB) based on output.

$$O(\frac{1}{\epsilon^4}|L|^3|N|^8\mathbb{L}\log\frac{|N|}{\epsilon} + |L|^3\mathbb{L}\log|N|)$$

- Efficiency enhancement:
 - \triangleright Approximate testing to shrink initial bound s.t. UB <= constant * LB.

$$O(\frac{1}{\epsilon^4}|L|^3|N|^4\mathbb{L}\log\frac{|N|}{\epsilon} + |L|^3|N|^4\mathbb{L}\log\log|N|)$$



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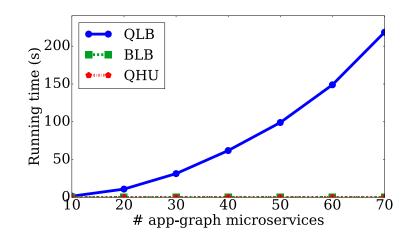


Simulation Settings

- Simulated network scenarios:
 - **❖** App-Graph:
 - > 5-layered applications, layer-1 being the source layer
 - > 10-70 microservices: 10% in layer-1, uniformly distributed in other layers
 - → 4 in-going edges per microservice in layers 2-5
 - Data distribution ratio: uniformly generated
 - ❖ Inf-Graph:
 - ➤ I instance per microservice in source layer, 5-15 in others
 - Linking probability (between interdependent instances): 0.3
 - Source demands: 100-900 units
 - ➤ Node capacities: 10-90 units
 - ➤ Node/Link delays: 0-500/1000 ms
 - Load balancing goal: optimal load under BLB, or 2 x optimal load under BLB
 - **\Leftrightarrow** Approximation parameter: ϵ =0.5
- Comparisons:
 - QLB
 - ♣ BLB
 - QHU: QoS-aware heuristic, solving BLB minimizing demand-weighted delay



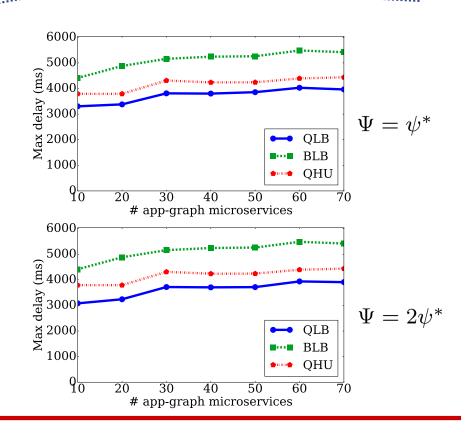
Comparison Results



Running time Polynomially increased

Max delay (QoS)

QLB has improved delay over the other solutions.





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Other Perspectives and Beyond

- So far, we've talked about
 - * Basic model: <u>DAG-based apps</u>, <u>Real-Graphs</u>
 - Processing capacities and delays
 - Network delays

Computing Perspective

- What we didn't consider in this work
 - Network topology
 - Network capacities & congestion
 - Routing
 - Reliability: microservice instance failures
 - Incentives, pricing
 - Payment methods

Networking Perspective

- Security Perspective

Economics Perspective

A unified approach is still in need for high-performance IoT.



Our Conclusions

- Load Balancing for Interdependent IoT Microservices
 - ❖ DAG model for applications: App-Graph and Inf-Graph
 - ❖ Application realization with Real-Graph abstraction
 - System-wide load balancing with QoS (delay) constraints
- Algorithmic solutions
 - Optimal solution for QoS-agnostic problem
 - FPTAS for (NP-hard) QoS-aware problem
- Future directions
 - Unified framework for IoT performance optimization



Thank you very much!

Q&A?



NP-Hardness Proof

